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Motion Planning: Overview and Foundations

CS 6301 Special Topics: Introduction to Robot Manipulation and Navigation

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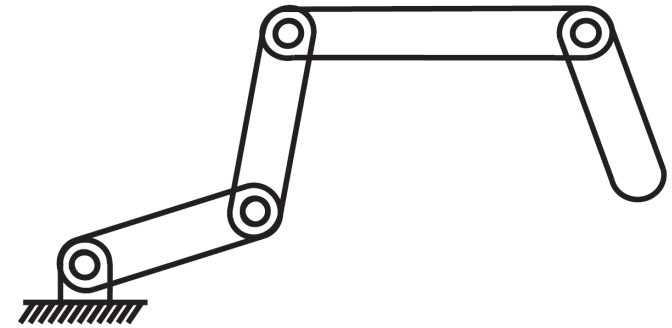
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Motion Planning

- Motion planning: finding a robot motion from a start state to a goal state (A to B)
 - Avoids obstacles
 - Satisfies other constraints such as joint limits or torque limits
- Recall configuration space (C-space)

Configuration Space

- The configuration of a robot is a complete specification of the position of every point of the robot.
- The minimum number n of real-valued coordinates needed to represent the configuration is the number of degrees of freedom (DOF) of the robot.
- The n -dimensional space containing all possible configurations of the robot is called the configuration space (C-space).
- The configuration of a robot is represented by a point in its C-space.



- 4 revolute joints
- 4 DOFs

Configuration Space

- The configuration of a robot arm with n joints
 - n joint positions $q = (\theta_1, \dots, \theta_n)$
- Free C-space $\mathcal{C}_{\text{free}}$
 - Configurations where the robot neither penetrates an obstacle nor violated a joint limit
- Robot state $x = (q, v) \in \mathcal{X}$
 - For second order dynamics, state is configuration and velocity $v = \dot{q}$
 - When using velocities as control input, state is the configuration $q(x)$
$$\mathcal{X}_{\text{free}} = \{x \mid q(x) \in \mathcal{C}_{\text{free}}\}$$

Equations of Motion

- The equations of motion of a robot

$$\dot{x} = f(x, u) \quad \text{Forward dynamics}$$

Robot state Control inputs $u \in \mathcal{U} \subset \mathbb{R}^m$

- Integral form

$$x(T) = x(0) + \int_0^T f(x(t), u(t)) dt$$

Motion Planning

- Given an initial state $x(0) = x_{\text{start}}$ and a desired final state x_{goal} find a time T and a set of control $u : [0, T] \rightarrow \mathcal{U}$ such that the motion

$$x(T) = x(0) + \int_0^T f(x(t), u(t)) dt$$

satisfies

$$x(T) = x_{\text{goal}}$$

$$q(x(t)) \in \mathcal{C}_{\text{free}} \text{ for all } t \in [0, T]$$

Types of Motion Planning Algorithms

- Path planning vs. motion planning
 - Path planning is a purely geometric problem of finding a collision-free path
$$q(s), s \in [0, 1] \quad q(0) = q_{\text{start}} \quad q(1) = q_{\text{goal}}$$
 - No concern about dynamics/control inputs
- Control inputs: $m = n$ **versus** $m < n$
 - When $m < n$, the robot cannot follow many paths
 - E.g., a car, $n = 3$, $m = 2$
- Online vs. Offline
 - Online is needed when the environment is dynamic

Types of Motion Planning Algorithms

- Optimal vs. satisficing

- In addition to reaching the goal state, we might want the motion planner to minimize a cost

$$J = \int_0^T L(x(t), u(t)) dt$$

Time-optimal $L=1$

Minimum-effort $L = u^T(t)u(t)$

- Exact vs. approximate

- Approximate $\|x(T) - x_{\text{goal}}\| < \epsilon$

- With or without obstacles

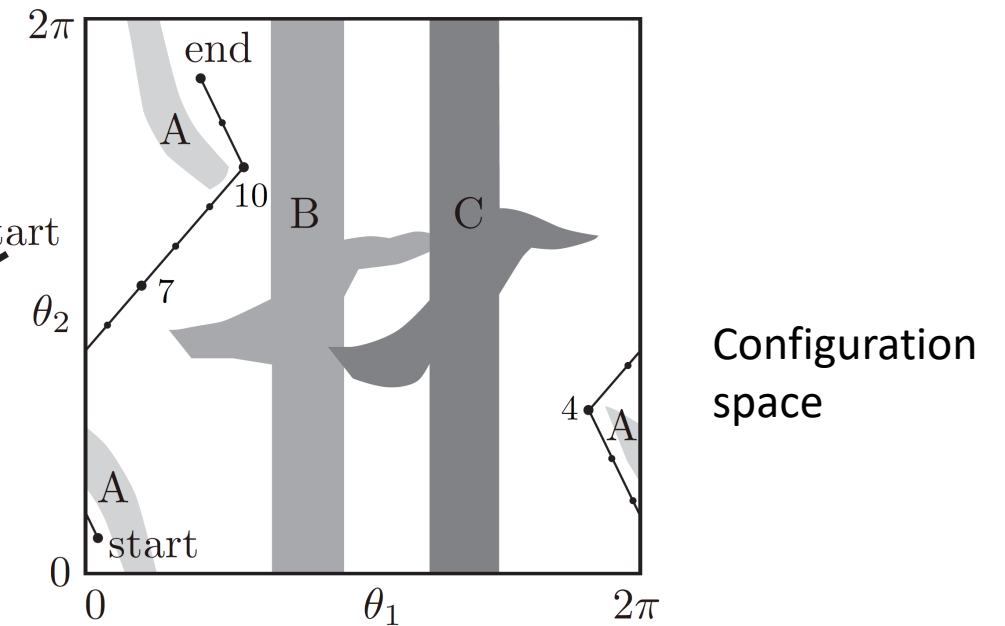
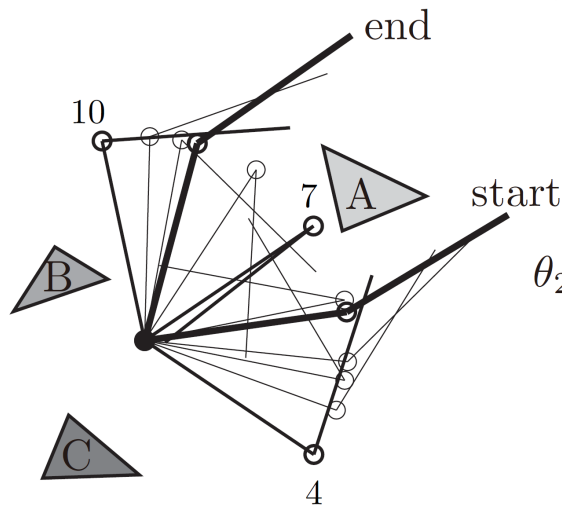
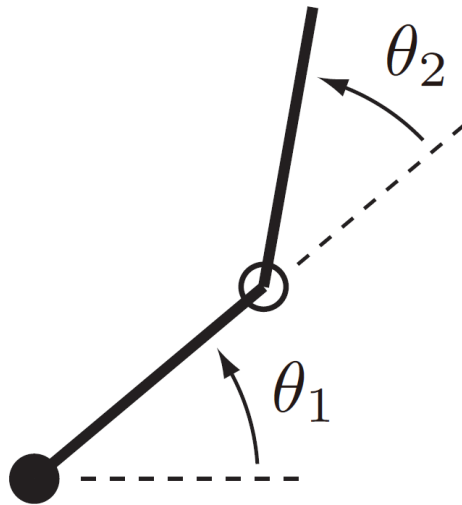
- Some motion planning problems are challenging even without obstacles
 - When $m < n$ or optimality is desired

Properties of Motion Planners

- Multiple-query vs. single-query planning
 - Multiple-query can build a data structure for $\mathcal{C}_{\text{free}}$
- “Anytime” planning
 - Continues to look for a better solution after a first solution is found
 - The planner can be stopped at anytime
- Completeness
 - A motion planner is said to be complete if it is guaranteed to find a solution in finite time if one exists, and to report failure if there is no feasible motion plan
- Computational complexity
 - The amount of time the planner takes to run or the amount of memory it requires

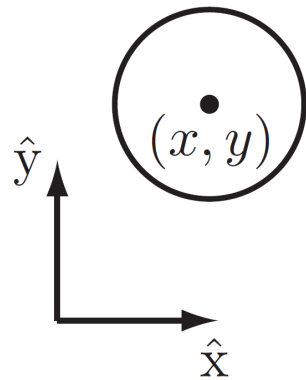
Configuration Space Obstacles

- Workspace obstacles partition the configuration space into two sets
 - Free space and obstacle space $\mathcal{C} = \mathcal{C}_{\text{free}} \cup \mathcal{C}_{\text{obs}}$
 - Joint limits are treated as obstacle in the configuration space
- A 2R planar arm

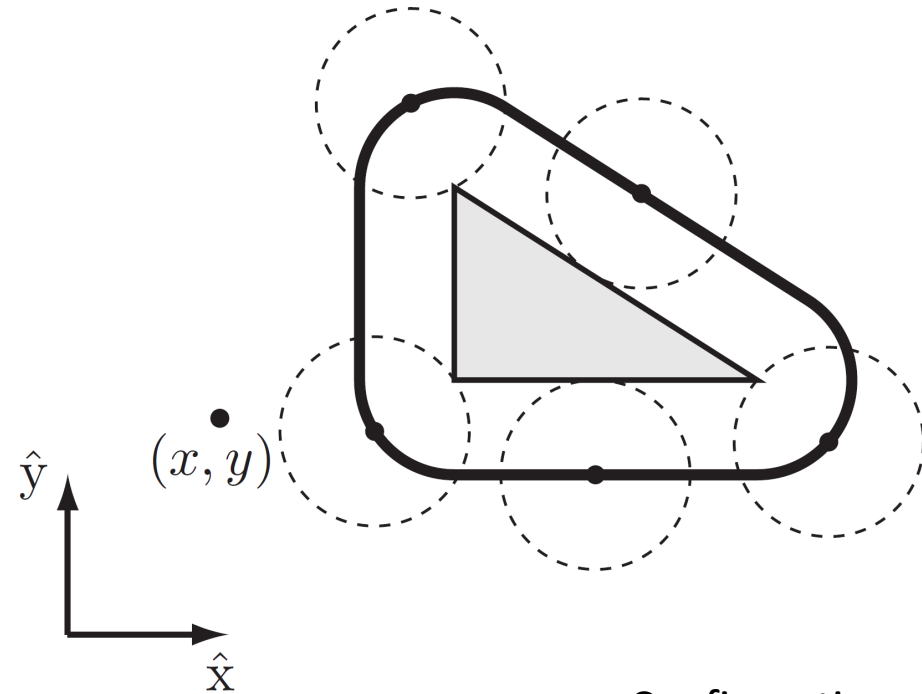
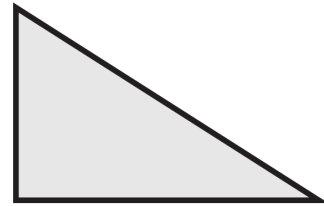


Configuration Space Obstacles

- A circular planar mobile robot



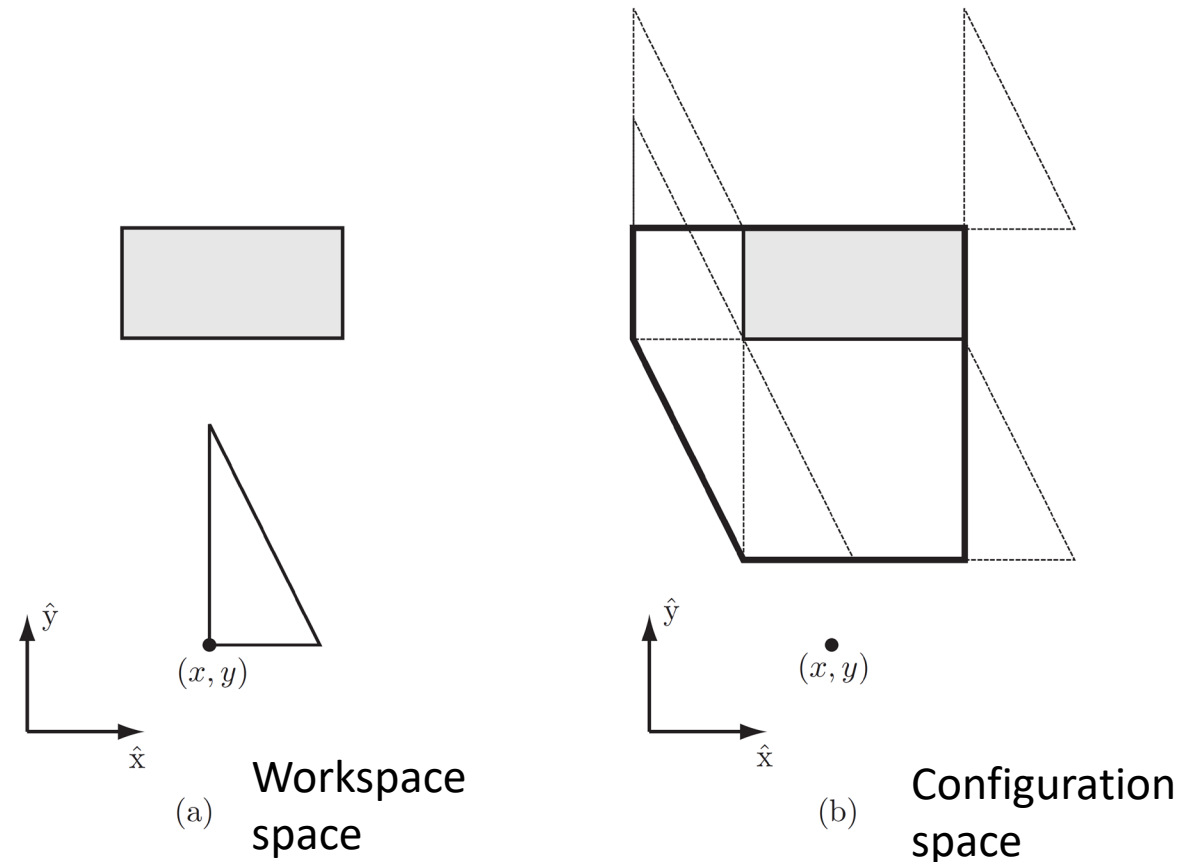
(a) Workspace space



(b) Configuration space

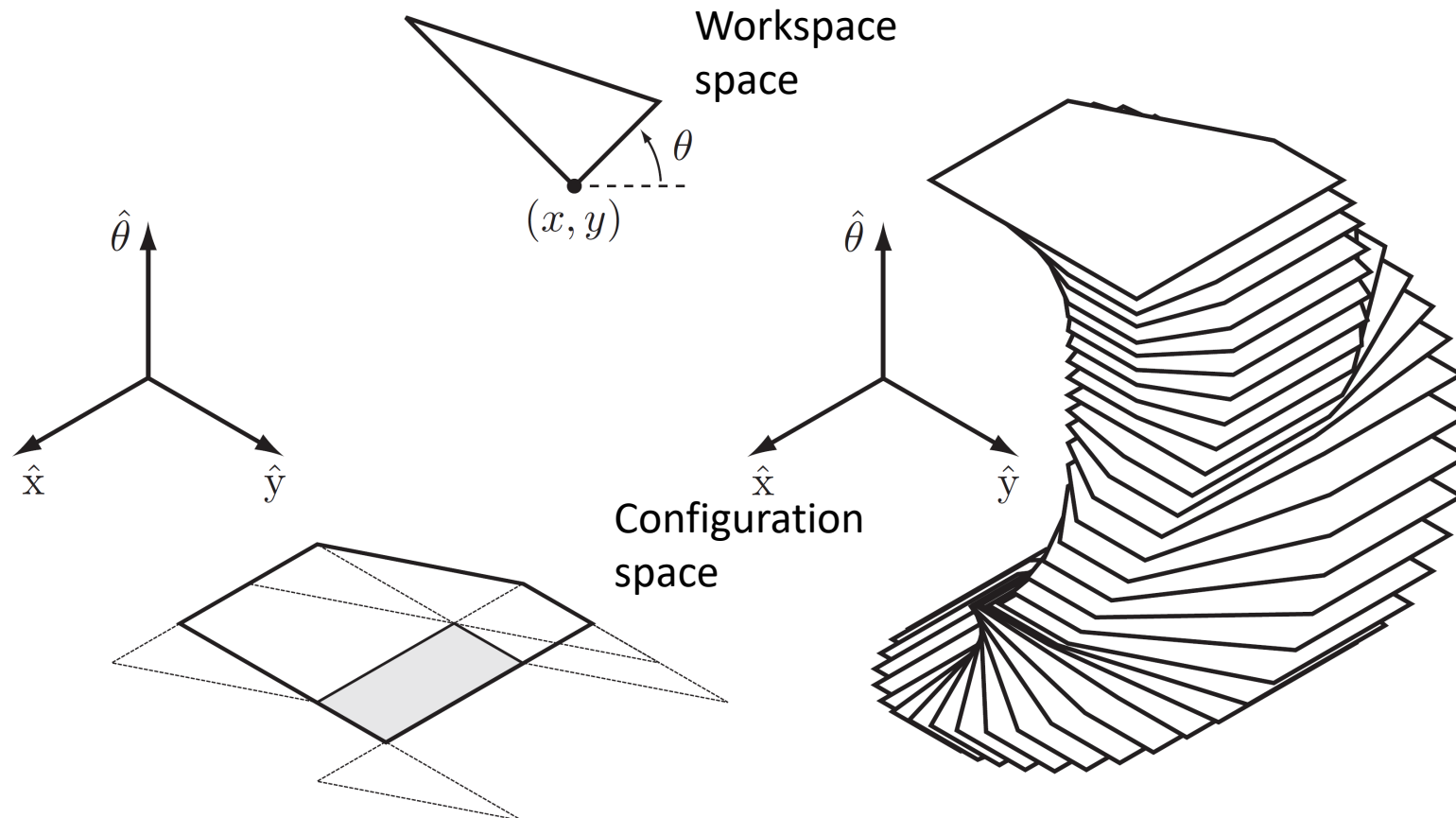
Configuration Space Obstacles

- A Polygonal Planar Mobile Robot That Translates



Configuration Space Obstacles

- A Polygonal Planar Mobile Robot That Translates and Rotates



Distance to Obstacles

- Given a C-obstacle \mathcal{B} and a configuration q , the distance between a robot and the obstacle

$d(q, \mathcal{B}) > 0$ (no contact with the obstacle),

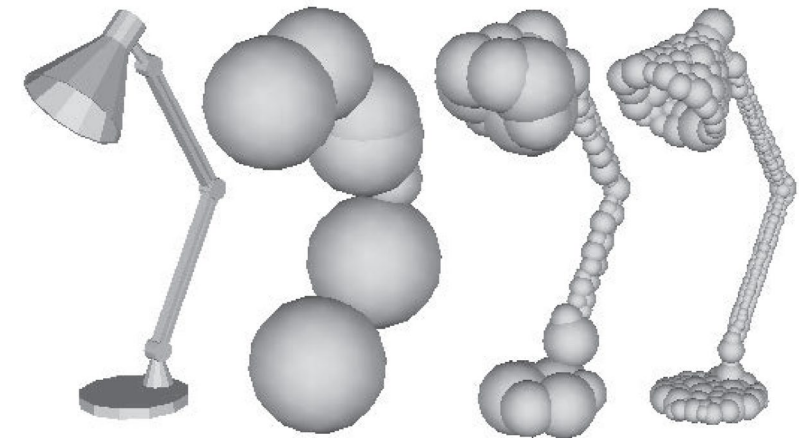
$d(q, \mathcal{B}) = 0$ (contact),

$d(q, \mathcal{B}) < 0$ (penetration).

- A distance measurement algorithm determines $d(q, \mathcal{B})$
- A collision detection algorithm determines whether $d(q, \mathcal{B}_i) \leq 0$

Distance to Obstacles

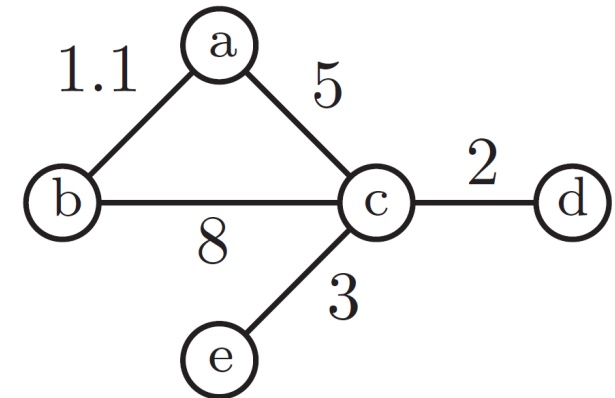
- Approximation of 3D shapes using 3D spheres
- Robot k spheres of radius R_i centered at $r_i(q)$
- Obstacle l spheres of radius B_j centered at b_j
- The distance between the robot and the obstacle



$$d(q, \mathcal{B}) = \min_{i,j} \|r_i(q) - b_j\| - R_i - B_j$$

A* Search Algorithm

- Finds a minimum-cost path on a graph
- Cost: sum of the positive edge costs along the path
- Data structures used
 - OPEN: a list of nodes not explored yet
 - CLOSE: a list of nodes explored already
 - $\text{cost}[\text{node1}, \text{node2}]$: positive, edge cost, negative, no edge
 - $\text{past_cost}[\text{node}]$: minimum cost found so far to reach node from the start node
 - $\text{parent}[\text{node}]$: a link to the node preceding it in the shortest path found so far



A* Search Algorithm

- Initialization
 - The matrix cost is constructed to encode the edges
 - OPEN is the start node 1
 - $\text{past_cost}[1] = 0$, $\text{past_cost}[\text{node}] = \text{infinity}$
- At each step
 - Remove the first node from OPEN and call it current
 - The node current is added to CLOSE
 - If current in the goal set, finished
 - Otherwise, for each neighbor of current that is not in CLOSE, compute

$$\begin{aligned} & \text{tentative_past_cost} \\ = & \text{past_cost}[\text{current}] + \text{cost}[\text{current}, \text{nbr}] \end{aligned}$$

A* Search Algorithm

- At each step (continued)
 - If $\text{tentative_past_cost} < \text{past_cost}[\text{nbr}]$
 $\text{past_cost}[\text{nbr}] = \text{tentative_past_cost}$
 $\text{parent}[\text{nbr}]$ is set to current
Compute estimated total cost for nbr
 $\text{est_total_cost}[\text{nbr}] \leftarrow \text{past_cost}[\text{nbr}] + \text{heuristic_cost_to_go}(\text{nbr})$
Add nbr to the correct position in OPEN (a sorted list)

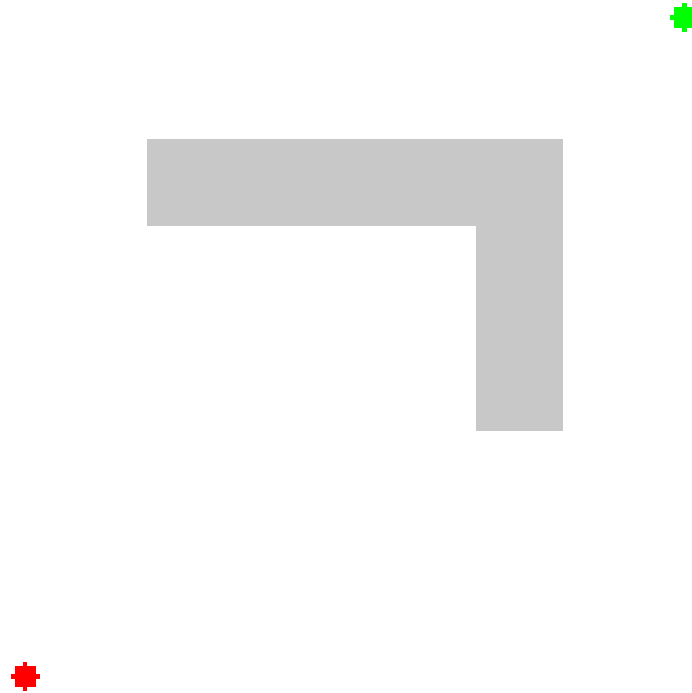
A* Search Algorithm

Algorithm 10.1 A* search.

```
1: OPEN  $\leftarrow$  {1}
2: past_cost[1]  $\leftarrow$  0, past_cost[node]  $\leftarrow$  infinity for node  $\in$  {2, ..., N}
3: while OPEN is not empty do
4:   current  $\leftarrow$  first node in OPEN, remove from OPEN
5:   add current to CLOSED
6:   if current is in the goal set then
7:     return SUCCESS and the path to current
8:   end if
9:   for each nbr of current not in CLOSED do
10:    tentative_past_cost  $\leftarrow$  past_cost[current] + cost[current, nbr]
11:    if tentative_past_cost < past_cost[nbr] then
12:      past_cost[nbr]  $\leftarrow$  tentative_past_cost
13:      parent[nbr]  $\leftarrow$  current
14:      put (or move) nbr in sorted list OPEN according to
          est_total_cost[nbr]  $\leftarrow$  past_cost[nbr] +
          heuristic_cost_to_go(nbr)
15:    end if
16:  end for
17: end while
18: return FAILURE
```

- Guaranteed to return a minimum-cost path
- Best-first searches

A* Search Algorithm



https://en.wikipedia.org/wiki/A*_search_algorithm

Summary

- Overview of motion planning
- Configuration space obstacle
- Distance to obstacles
- A* search algorithm

Further Reading

- Chapter 10 in Kevin M. Lynch and Frank C. Park. Modern Robotics: Mechanics, Planning, and Control. 1st Edition, 2017.
- A* search: P. E. Hart, N. J. Nilsson, and B. Raphael. A formal basis for the heuristic determination of minimum cost paths. IEEE Transactions on Systems Science and Cybernetics, 4(2):100-107, July 1968.